

Photon: Shining a light on the I/O problem in HPC

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The Problem

- ☀ Growing I/O gap in HPC
 - * Problems effectively using the network in Clusters, Clouds and Grids
- ☀ On the one hand -- Extreme scale systems
- ☀ On the other -- 100Gb/s networks
- ☀ Components scaling up but the gap is growing
 - * Highlighted in various places including yesterday's panel

One scenario

- ☀ MPI-IO to a parallel filesystem
- ☀ GridFTP servers mount this filesystem and perform parallel file transfers
- ☀ Data has been forced into a sequential file
- ☀ Did the parallelism in the program match that of the object stores in the filesystem?
Does the GridFTP striping match it?
- ☀ “Optimized” separately (if at all)
 - * Even separately, the optimizations are inadequate

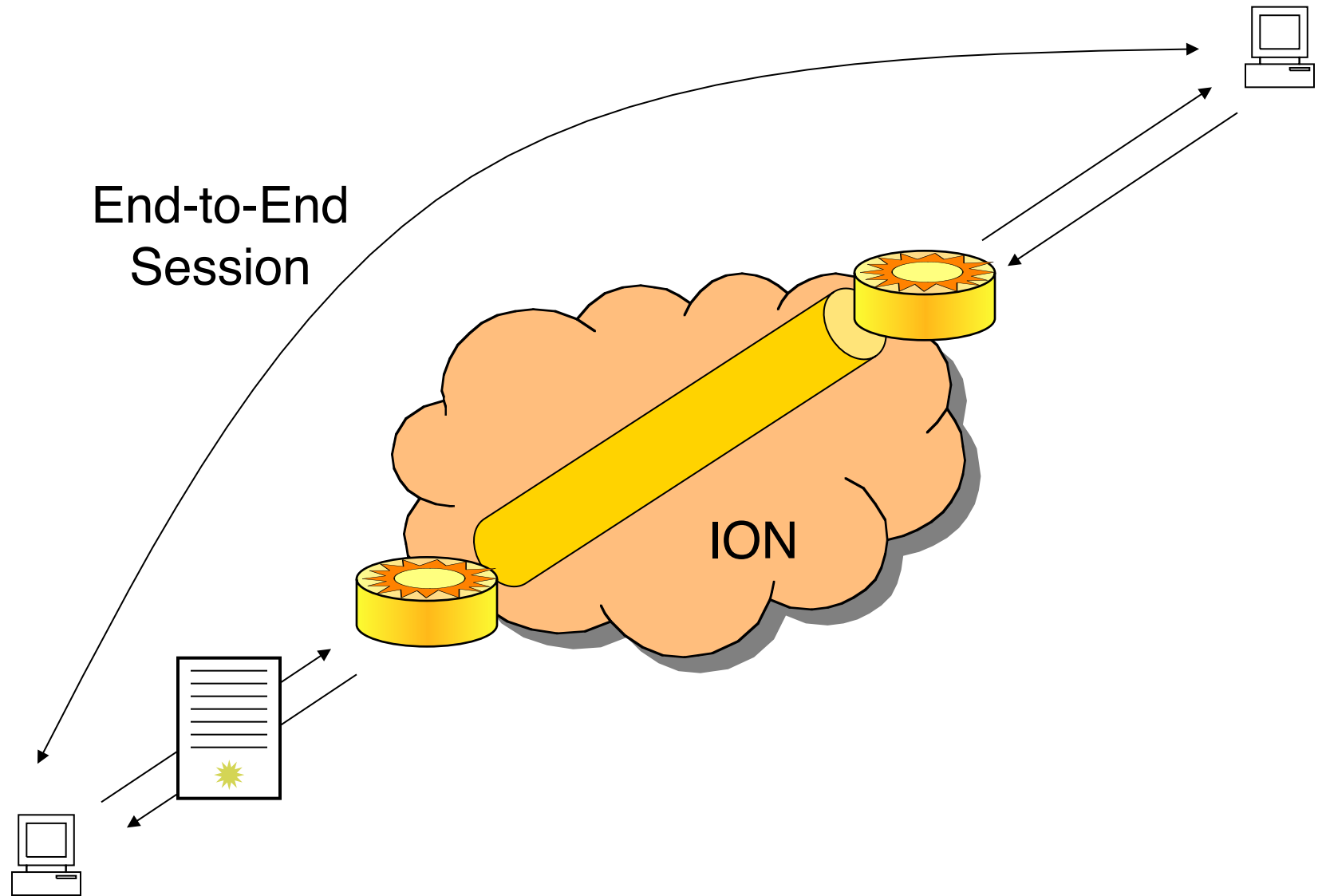
Photon

- ☀ Unifies previous solutions into an end to end system for parallel I/O
- ☀ Wide-area data movement
 - * Phoebus
- ☀ MPI program transformation
 - * AToMS
- ☀ Lightweight cluster data movement
 - * Gravel -> Photon
- ☀ Deconstructed filesystems
 - * eXnode

Ph☀ebus

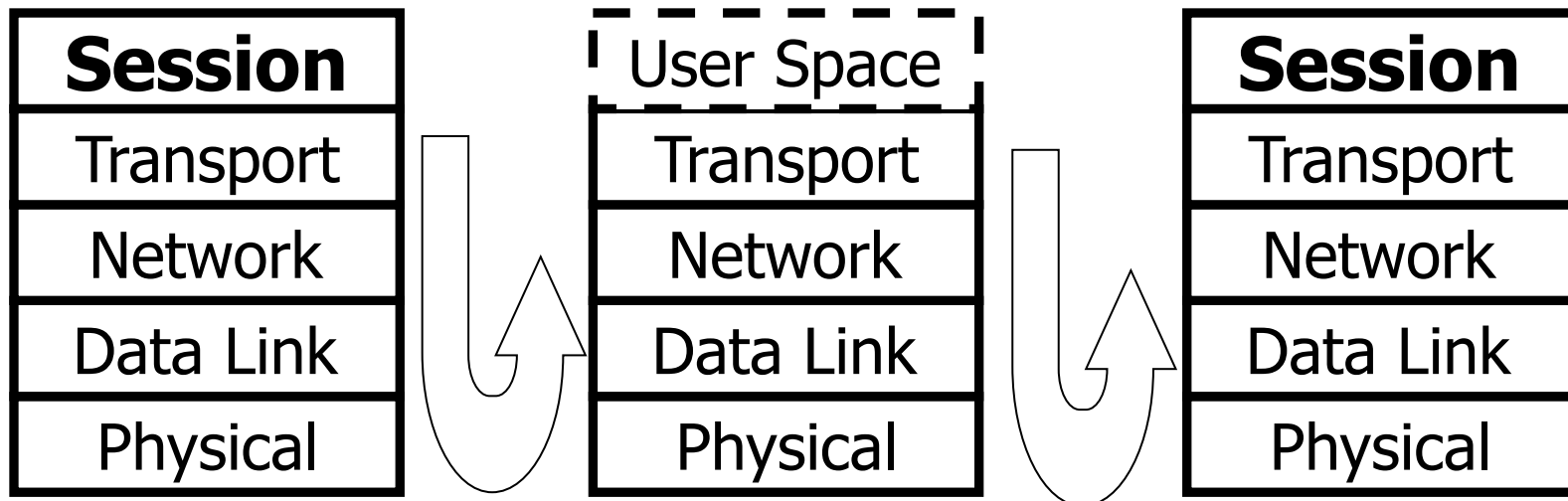
- ☀ The Phoebus project aims to bridge the network performance gap by providing an optimizing network service
- ☀ Phoebus is based on the concept of a “session” that enables multiple adaptation points in the network to be composed
- ☀ Phoebus provides a gateway for legacy applications to use advanced networks
 - * Network reservation like ESnet’s OSCARS, Internet2 ION

Ph☀ebus



Session Layer

- ☀ A *session* is the end-to-end composition of *segment-specific* transports and signaling
 - ✦ More responsive control loop via reduction of signaling latency
 - ✦ Adapt to local conditions with greater specificity
 - ✦ Buffering in the network means retransmissions need not come from the source



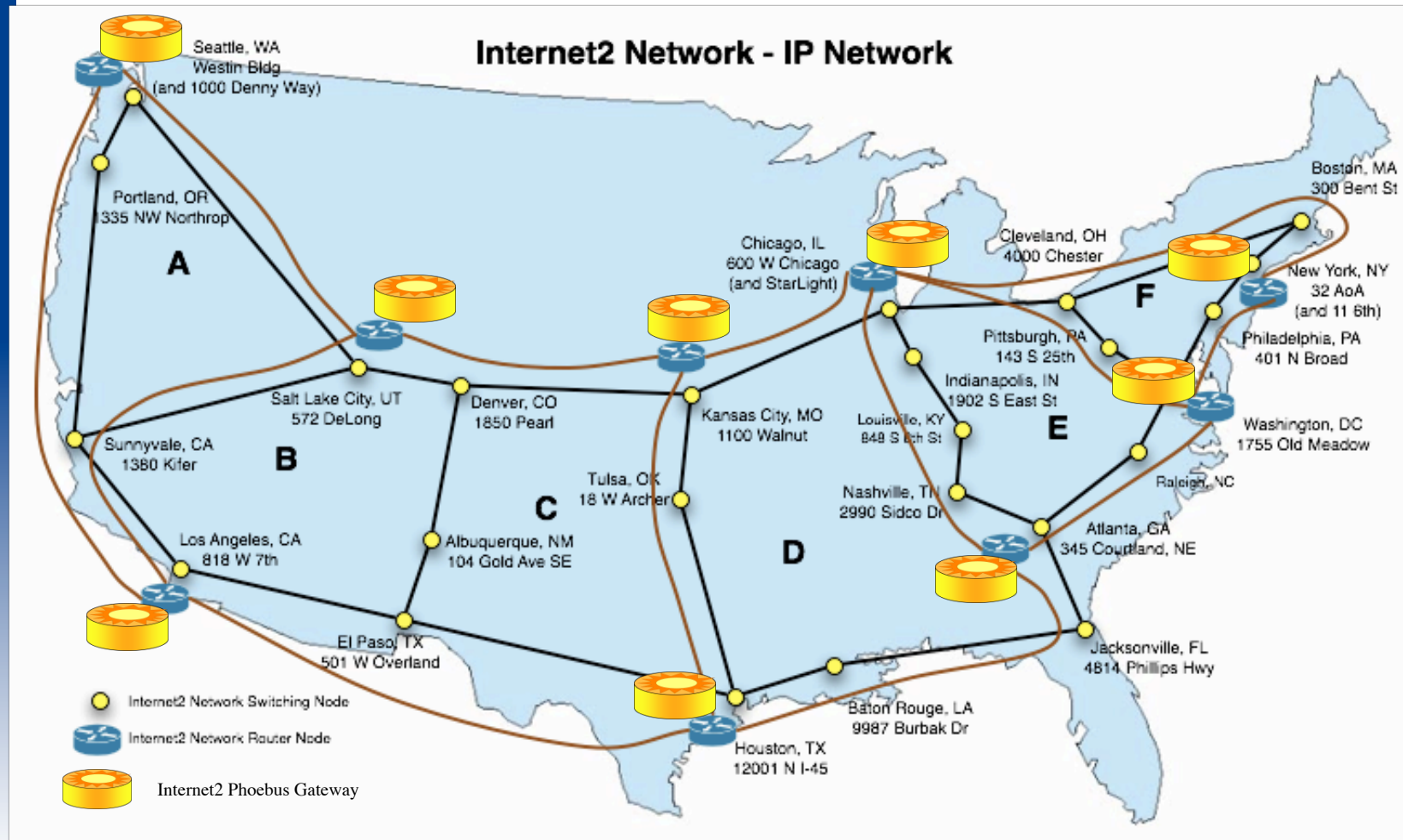
Dynamic Networks

- ☀ The last piece of the Cloud puzzle
 - ✦ Network allocation is the elephant in the cloud
- ☀ Phoebus signals dynamic networks like ESnet, Internet2, GEANT...
 - ✦ Phoebus speaks to the control plane to provision network resources
- ☀ Once the connection is established to the Phoebus node, traffic can begin to flow
 - ✦ Could be sent over an existing link if unable to provision
 - ✦ Phoebus can finish the connection over the commodity network if the allocation times out

Session Layer Benefits

- ☀ Our session-layer approach is an architectural evolution for the Internet
 - * ongoing work in GENI, DOE Networks program
- ☀ A session layer provides explicit control over ***adaptation points*** in the network
 - * Transport protocol
 - * Rate-based to congestion based
 - * Shorter feedback loops
 - * Traffic engineering
 - * Map between provider-specific DiffServ Code Points / VLANs
 - * Authorization and Authentication
 - * Rich expression of policy via e.g. the Security Assertion Markup Language (SAML)

Deployment Plans



Phoebus and GridFTP

- ☀ Integrated with GridFTP
 - * *globus-gridftp-server* loads the Phoebus XIO driver when requested
 - * *globus-url-copy* extended to support Phoebus-based transfers with -ph flag or explicitly with -dcstack
 - * Support for advanced features
 - * 3rd party transfers
 - * Parallel streams
- ☀ SC09 paper and JPDC article in press show performance over 10G

AT MS

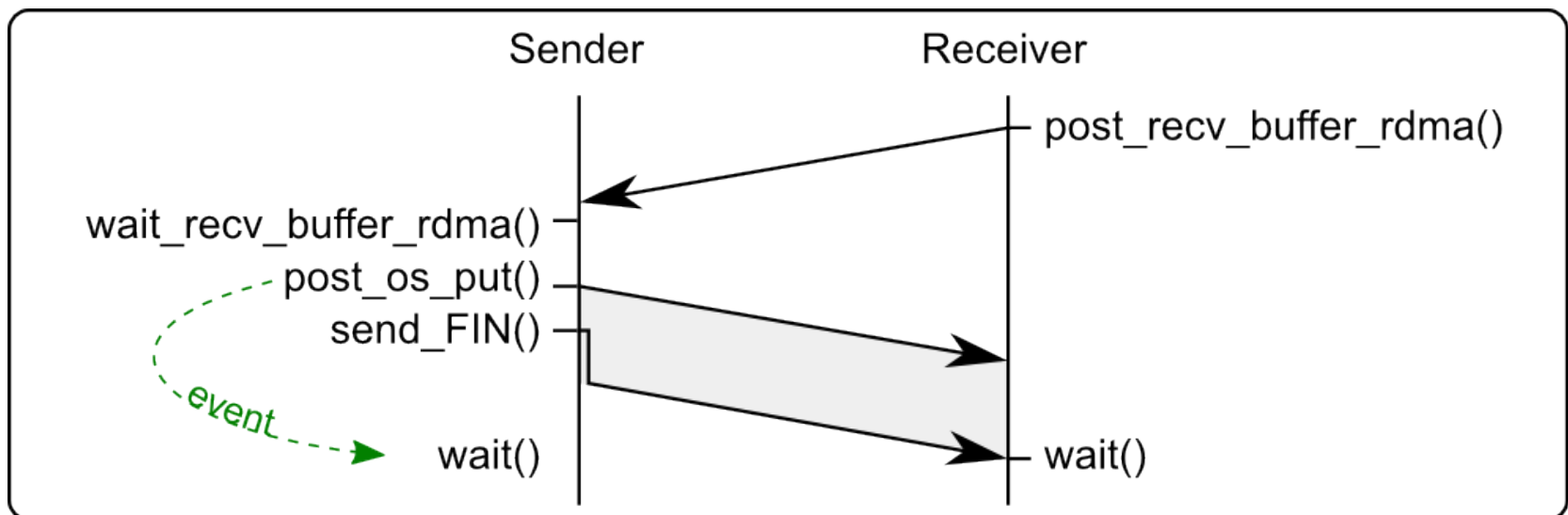
- ☀ Focused on message-passing performance in clusters
- ☀ AToMS = Auto-Tuning of MPI Software
 - * Actually, Transformation and Tuning
- ☀ Kennedy's telescoping languages work – “improves performance by replacing sequences of library calls with equivalent, but more efficient, sequences.”
- ☀ Partial implementations in Open64 and LLVM

AToMS Transformations

- ☀ Transform MPI communication
 - * Collectives → Point-to-point
 - * Blocking → Non-blocking
 - * Non-blocking → One-sided
 - * Send fission and fusion
 - * Restructure (user-defined) MPI datatypes
- ☀ Separate components of communication, code motion to improve overlap
 - * Memory registration
 - * Metadata exchange
 - * Data movement
 - * Progress/Completion

Gravel

- ☀ Library for use by AToMS transformations
- ☀ RDMA put/get for data movement
 - ✱ Also MX
- ☀ “Ledger” for progress and synchronization
 - ✱ Also available via RDMA

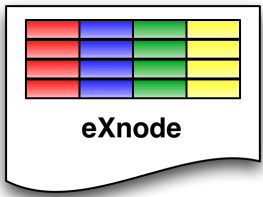
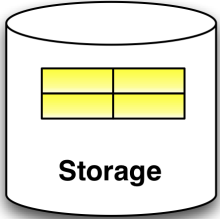
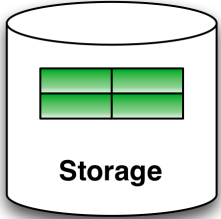
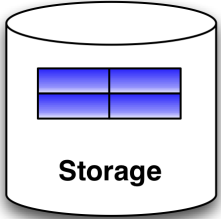
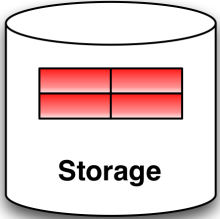
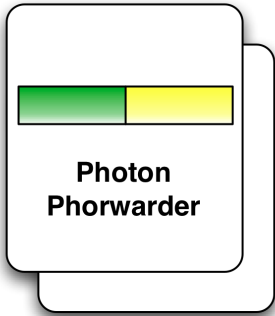
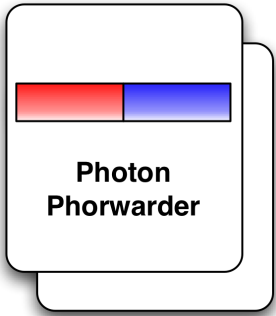
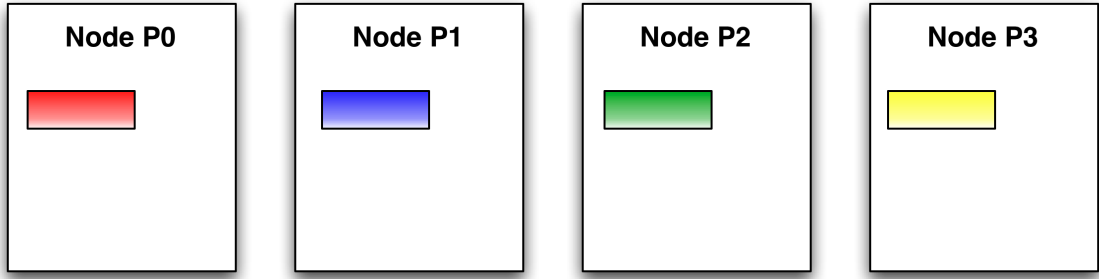


eXnode

- ☀ Concept from Logistical Networking work (IBP)
- ☀ Analogous to a filesystem inode, but available in the application
- ☀ Describes location and relationship of blocks in a (virtual) file
- ☀ Allows us to provide filesystem-like semantics without filesystem overhead
- ☀ Data “chunks” in the eXnode can refer to IBP allocations, Photon buffers, etc

Photon

- ☀ Transform MPI-IO calls into Photon calls
 - * pattern, location are often discoverable and consistent over the lifetime of the application
- ☀ Create data movement session with various intermediate forwarders
 - * Similar ideas to those in DART, ZOID
- ☀ This takes advantage of our session protocol and asynchronous progress notification
 - * “For the next 100K iterations, watch the ledger for completion, grab the data, update the completion ledger”
- ☀ Unify wide-area and cluster-area optimizations
 - * Building on the mature Phoebus forwarder



Implications for next-generation systems

- ☀ Programs need APIs, with flexibility of implementation
- ☀ In particular, when library or OS functionality can be inlined in the application, we expose more opportunities to mitigate latency
- ☀ Don't force a given application to pay for what it doesn't need – “deconstruction”
 - * E.g., POSIX I/O with filesystem semantics
- ☀ Single framework enables optimizations

End

☀ Thank you for your attention

☀ Questions?